



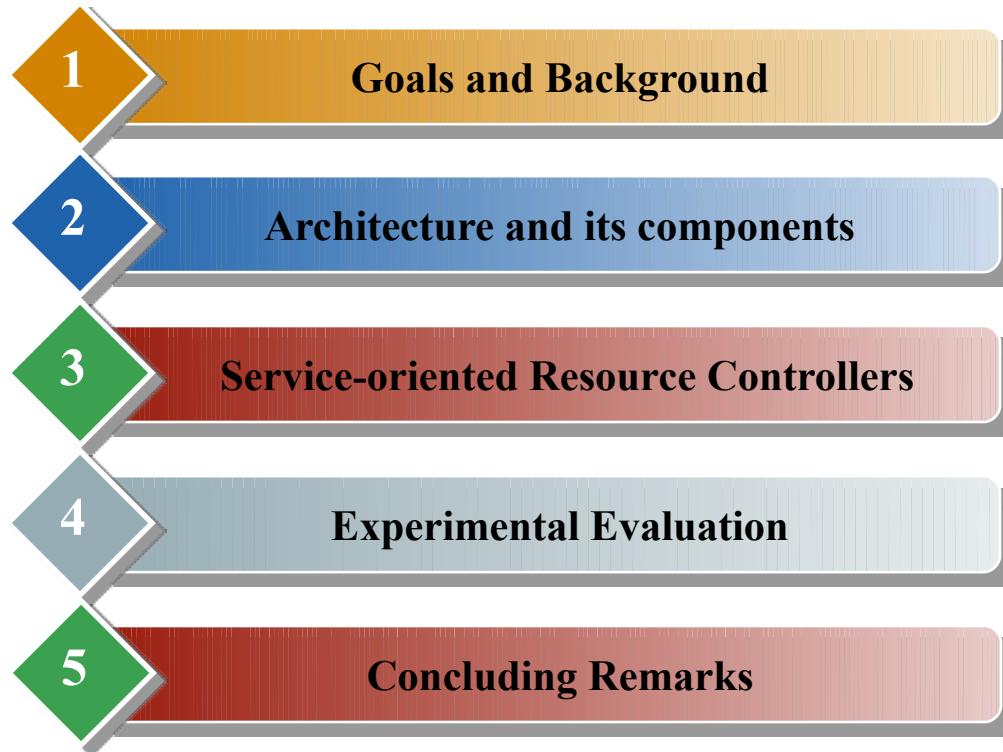
System and Virtualization Management : Standards and the Cloud

Corset :
Service-oriented Resource Management System in Linux

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Contents



Goal & Problems

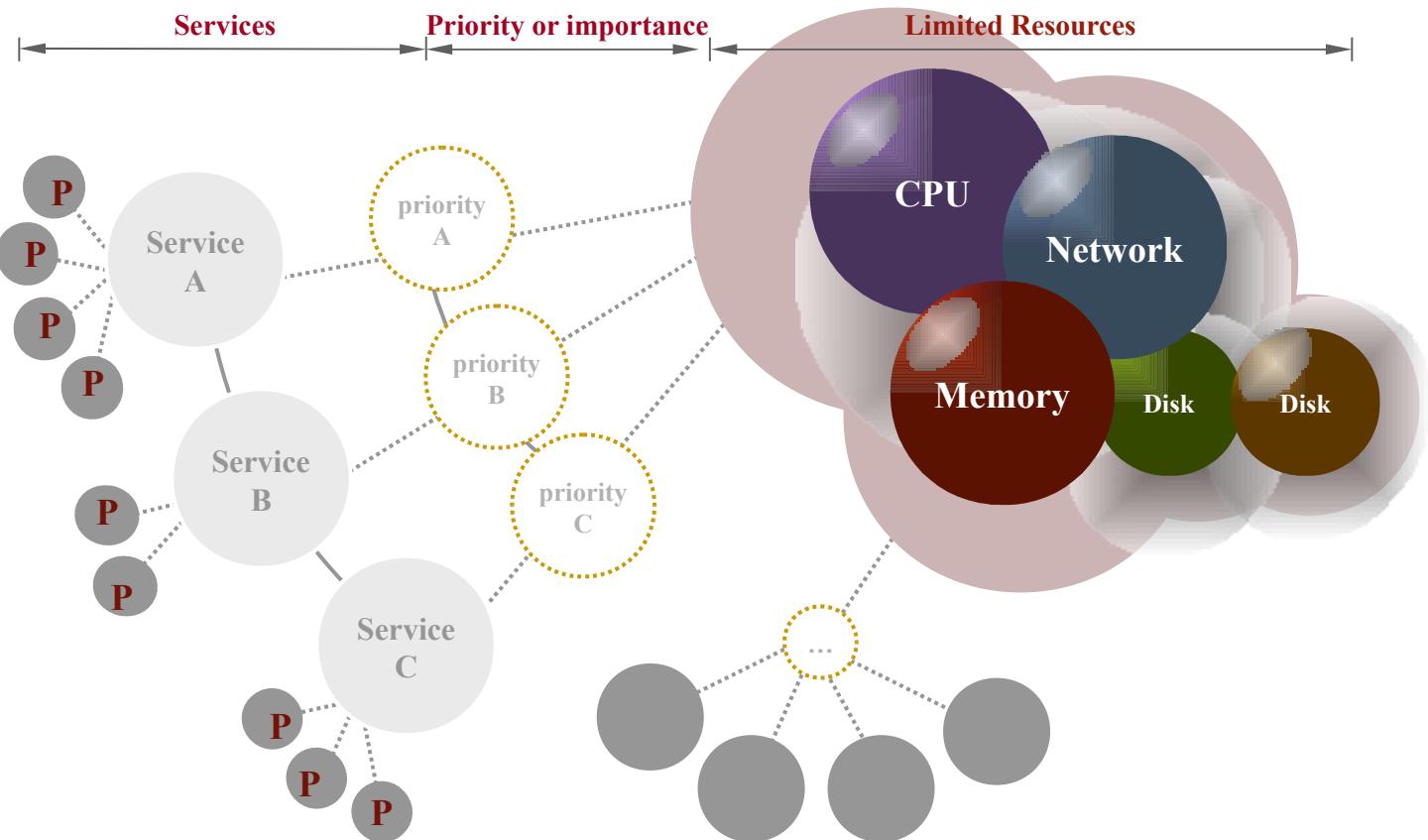
❖ Goal

- Guaranteeing the service QoS or stability in unexpected workload situation

❖ Problems

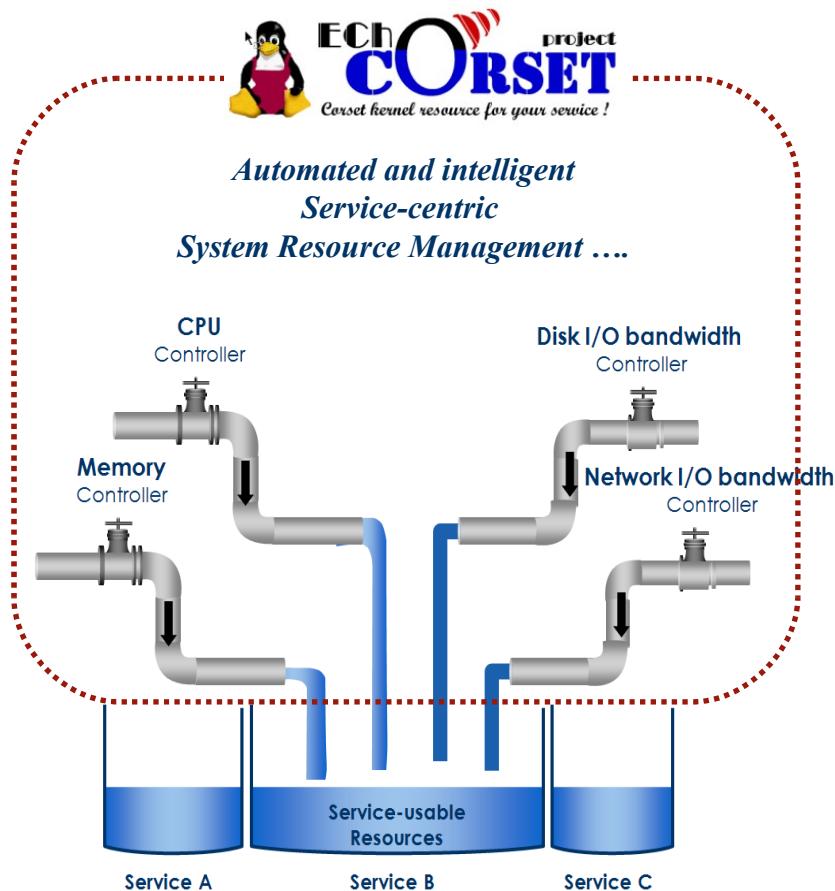
- System resources are not enough for many running services and application in a system
- As the service is more complex, process-unit or system-unit resource management is not enough for guaranteeing the service performance or QoS
- All services running in a system don't have same importance
- Administrator can't guarantee that the specific service has proper resources in unsettled workload situation
- System resource management is not easy for administrators

Background

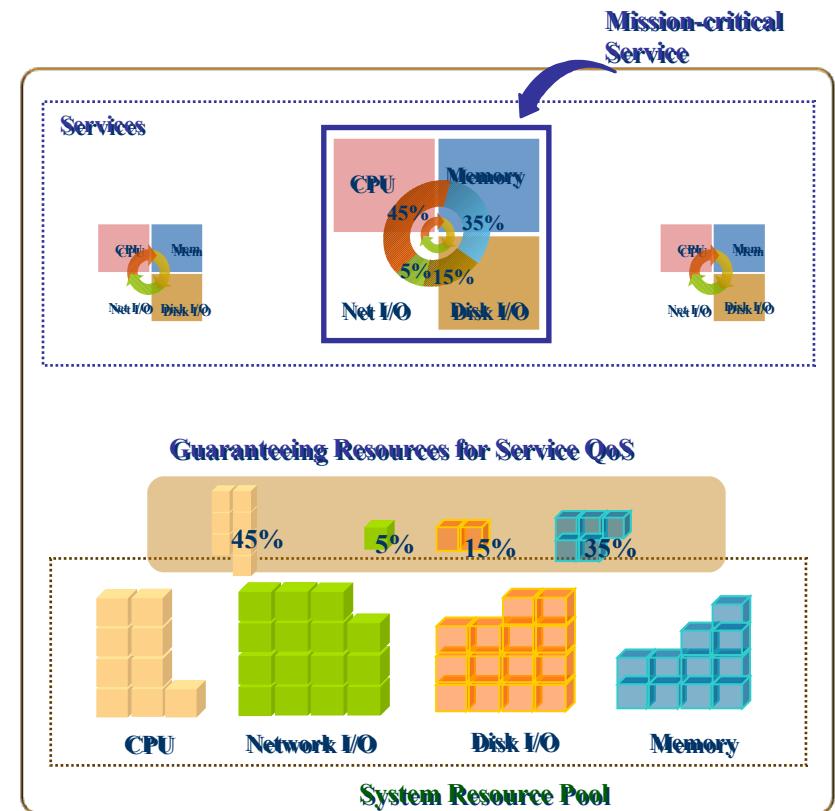


Concept of Corset

❖ What is?



❖ What can we do with?



Corset Architecture

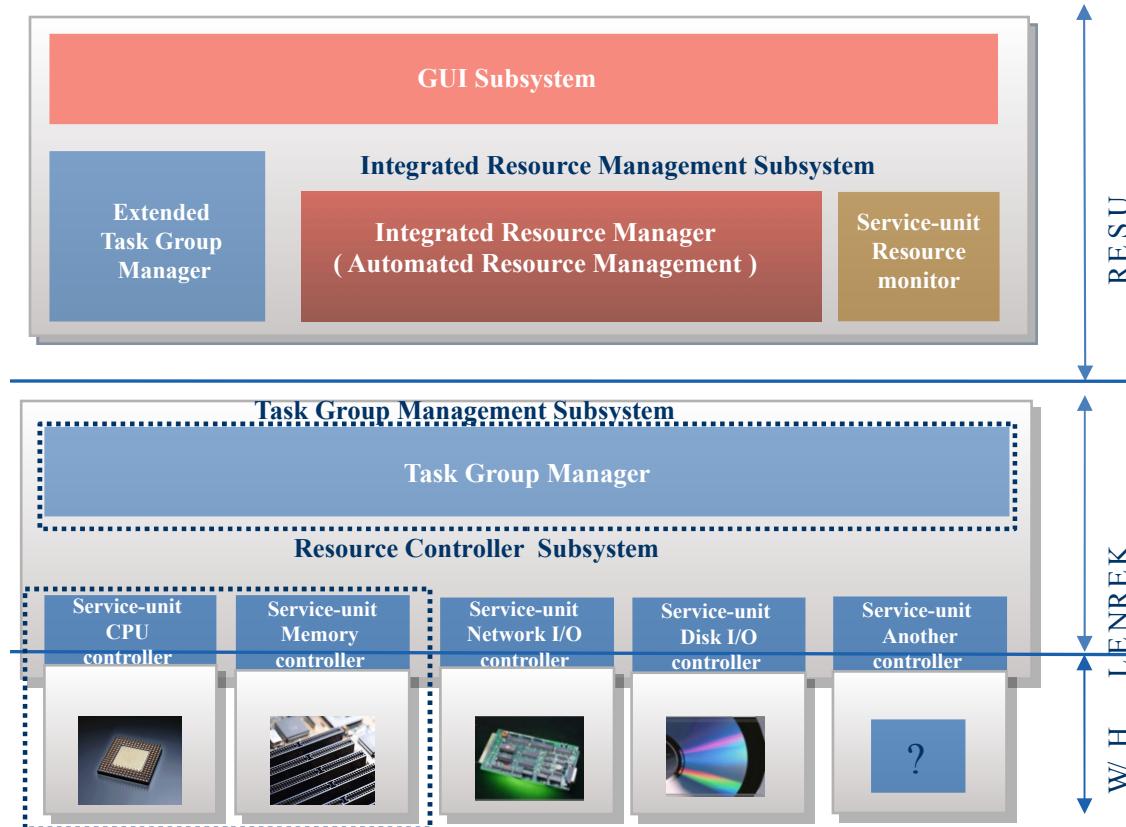


Fig. This is whole architecture of CORSET that consists of four subsystems.

Components of Corset (1/2)

❖ Resource controller subsystem

- May be a kind of resource scheduler or controller
 - Allocate or withdraw each resource to/from the specific service
 - Support predictable or proportional sharing of resources according to service priority
 - CPU, Memory, Disk I/O bandwidth, Network I/O bandwidth and so on

❖ Task group management subsystem

- Allow administrator to organize a new service with the processes he/she want to include based on PID
 - Create or destroy service
 - Add/delete/move processes to/from specific service

Components of Corset (2/2)

❖ Integrated management subsystem

- Support dynamic & autonomic reconfiguration according to changeable and unexpected workload
- create and provides service-unit resource information to assist service-unit resource management

❖ GUI management subsystem

- Support convenient and easy environment for system resource management by abstracting the complex interfaces of below blocks

Service-oriented Resource Controllers(1/3)

❖ Service-oriented CPU and Memory Controller

- We adapt existing functionalities in linux kernel (about upper kernel-2.6.24)
- CPU
 - CFS scheduler + cgroup framework
 - Proportional share
- Memory
 - mem_cgroup + page_cgroup
 - Limiting the maximum usage

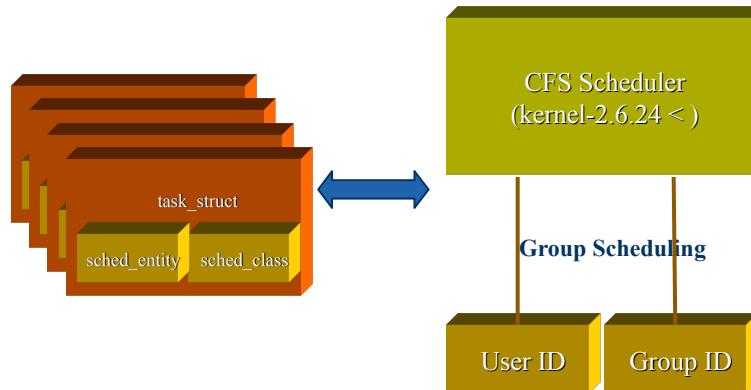


Fig. CPU Control using existing functionality

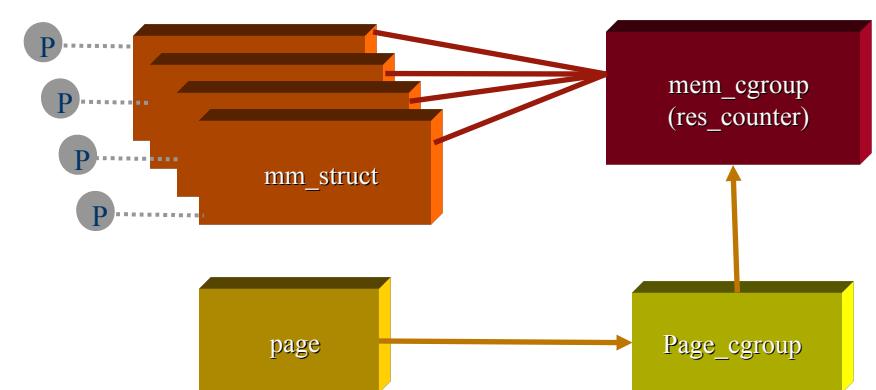


Fig. Memory Control using existing functionality

Service-oriented Resource Controllers(2/3)

❖ Service-oriented Disk I/O Controller

- A kind of new device mapper driver
- Independent of underlying specific I/O schedulers
- Supported policies
 - Proportional share(weight)
 - Range share(range-bw)

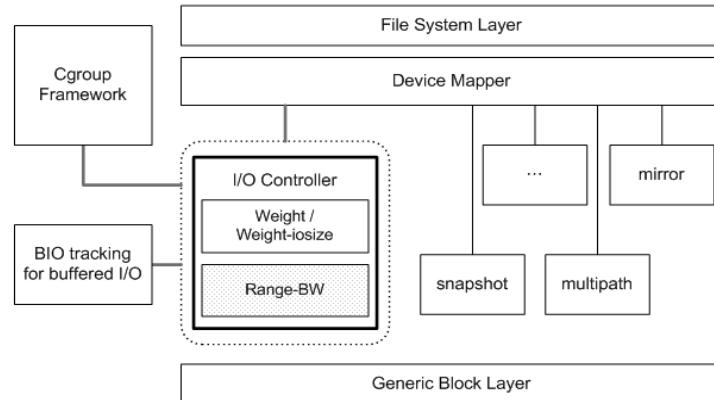


Fig. Overview of Disk I/O Controller

- <http://sourceforge.net/projects/ioband>

Service-oriented Resource Controllers(3/3)

❖ Service-oriented Network Controller

- TC(traffic control) module already support the limiting or fixed bandwidth
 - Based on IP, port and something like that
- TC didn't support the process group based bandwidth control
 - We added assigning, exporting and importing P.G ID in TC

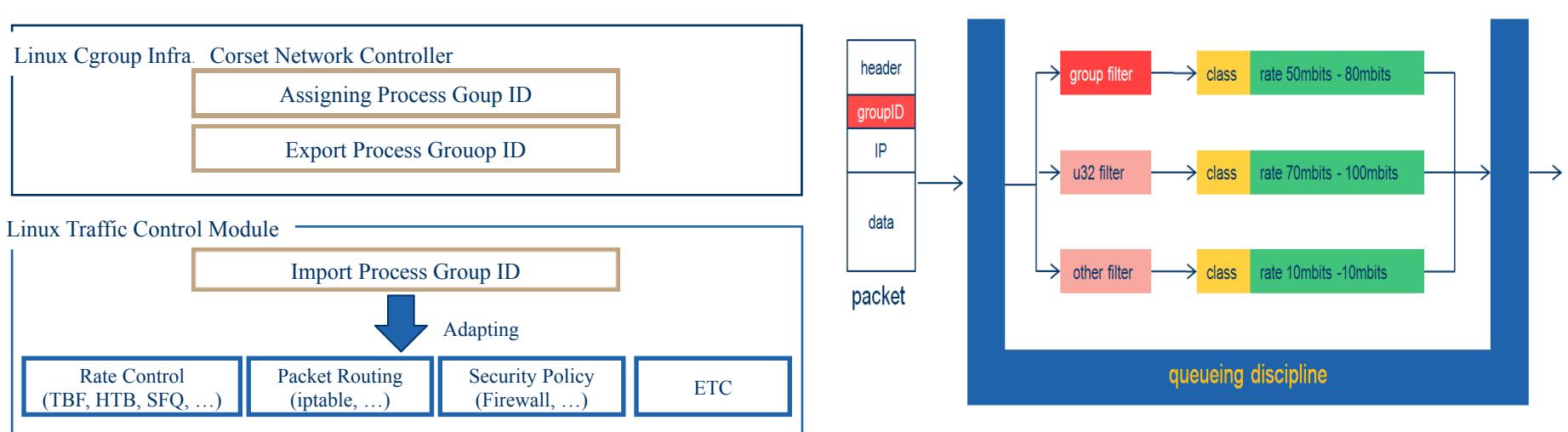


Fig. Overview of Netwrok I/O Controller

Evaluation Environment

❖ Category for Evaluation

- Resource guaranteeing test
 - groupA : groupB : groupC = 1 process : 50 or 100 processes : 50 or 100 processes
- Resource limitation test
 - groupA : groupB : groupC = 1 process : 1 processes : 50 or 100 processes
- Resource isolation test
 - groupA : groupB : groupC = 50 or 100 processes : 50 or 100 processes : 50 or 100 processes

❖ H/W and S/W Specification for test

Evaluation System Specification (H/W & S/W)						
	System	OS	Disk	CPU	Mem	Workloader
Spec	SAMSUNG Smart Server ZSS108	linux (kernel-2.6.30-rc1)	SAMSUNG SATA-2, 7,200rpm, 80G	2-way, Intel Pentium4 3.4 GHz	2 GB	CPU : ebizzy Mem : ebizzy Disk : fio-1.26 Net : iperf-2.0.4

Experimental Evaluation (Disk I/O BW)

Group1 : group2 : group3 = 2(1) : 1(100) : 1(100)

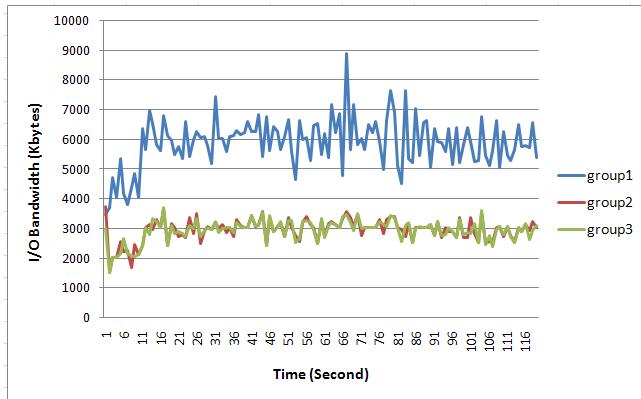


Fig. I/O bandwidth in resource guaranteeing test(proportional)

Group1 : group2 : group3 = 11~12M(1) : 5~6M(100) : 5~6M(100)

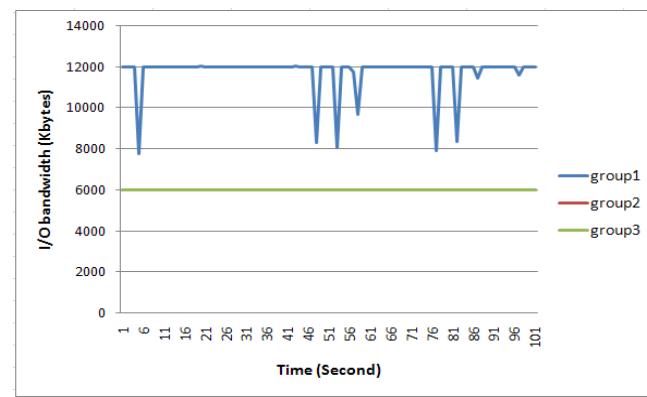


Fig. I/O bandwidth in resource guaranteeing test(range-bw)

Group1 : group2 : group3 = 1(100) : 2(1) : 2(1)

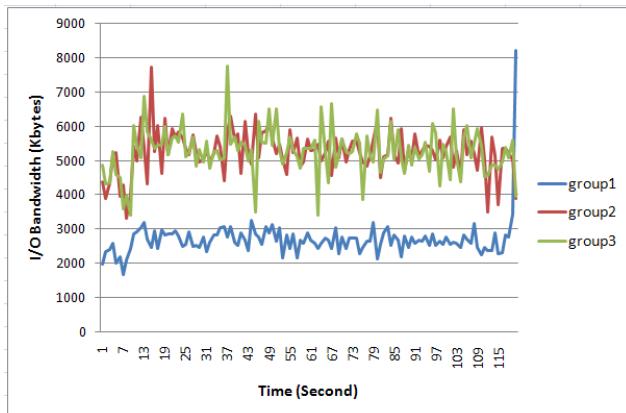


Fig. I/O bandwidth in resource limitation test(proportional)

Group1 : group2 : group3 = 1(100):5(100):9(100)

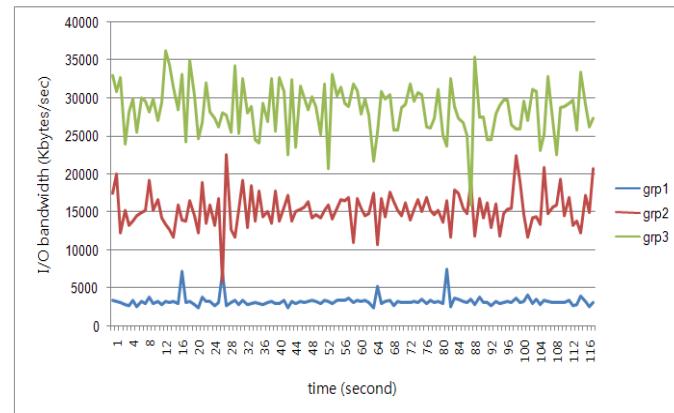


Fig. I/O bandwidth in resource isolation test(proportional)

Experimental Evaluation (Network I/O BW)

Group1 : group2 = 600Mbps(10) : 300Mbps(up to 100)

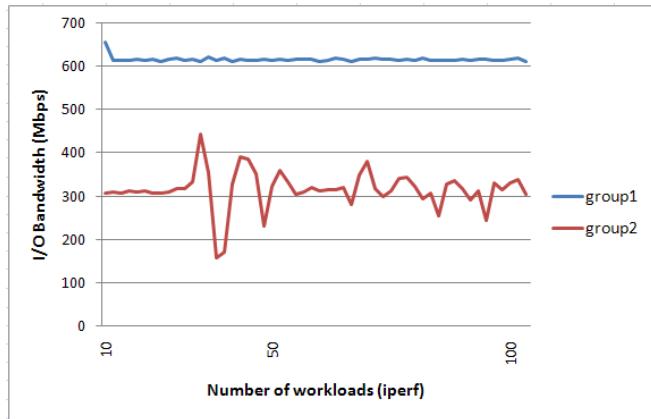


Fig. I/O bandwidth in resource guaranteeing test

Group1 : group2 = 700Mbps(10) : 200Mbps(up to 100)

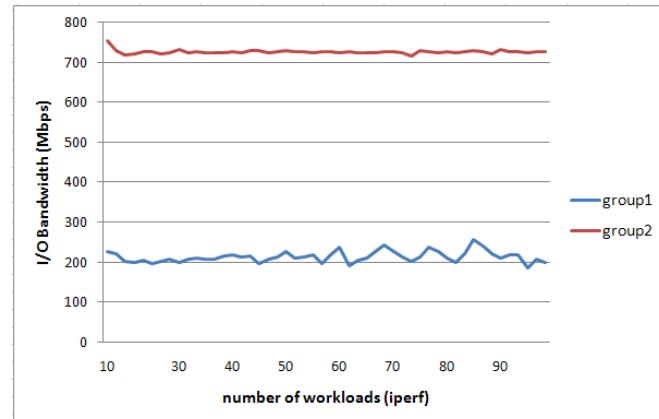


Fig. I/O bandwidth in resource limitation test

Group1 : group2 : group3 = 10Mbps(50) : 30Mbps(50) : 60Mbps(50)

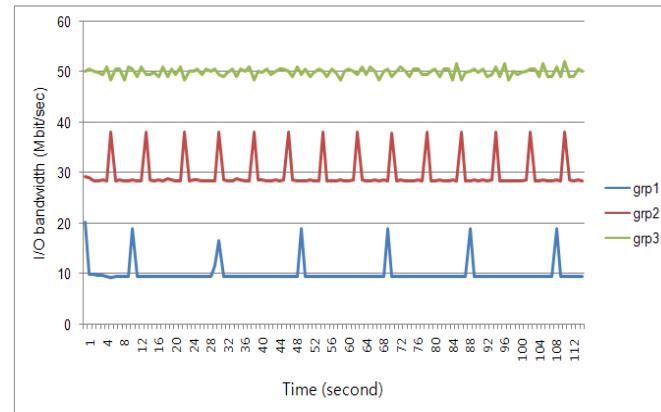


Fig. I/O bandwidth in resource isolation test

Experimental Evaluation (Memory)

Group1 : group2 : group3= 10M(1) : 200M(1) : 700M(1)

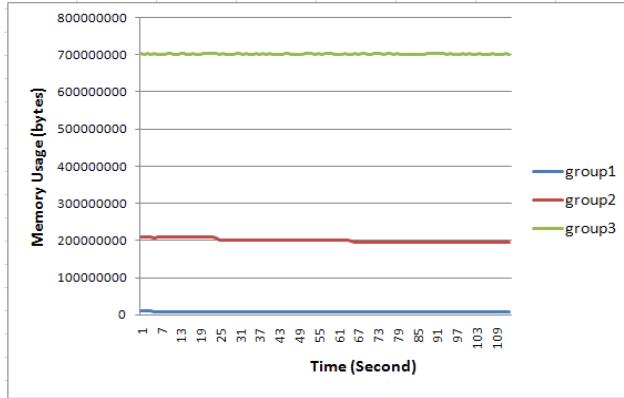


Fig. Memory control in normal test

Group1 : group2 : group3= 500M(1) : 250M(50) : 250M(50)

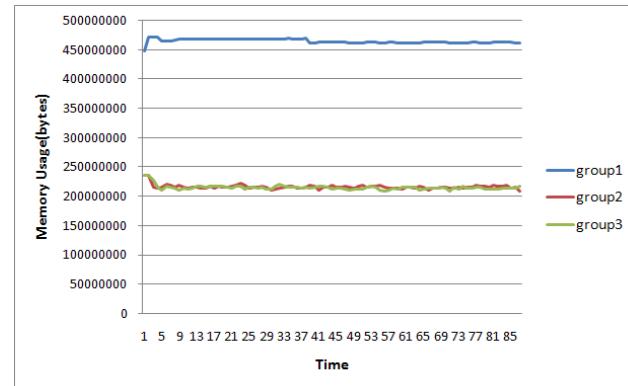


Fig. Memory control in resource guaranteeing test

Group1 : group2 : group3= 200M(50) : 400M(1) : 400M(1)

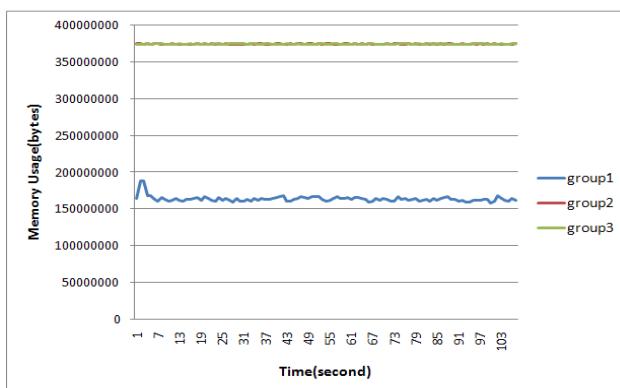


Fig. Memory control in resource limitation test

Group1 : group2 : group3= 500M(50) : 333M(50) : 166M(50)

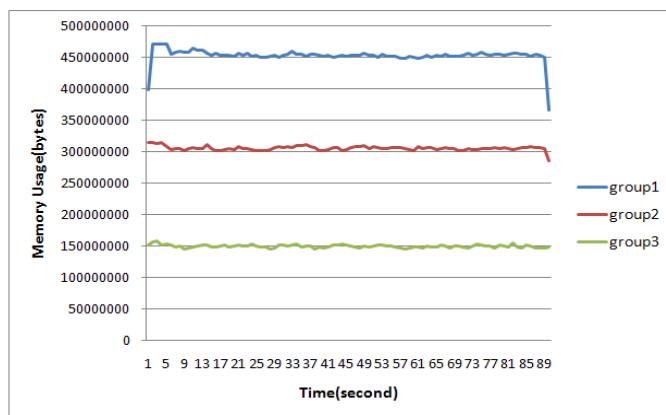


Fig. Memory control in resource isolation test

Experimental Evaluation (CPU)

Group1 : group2 = 2(1) : 1(200)

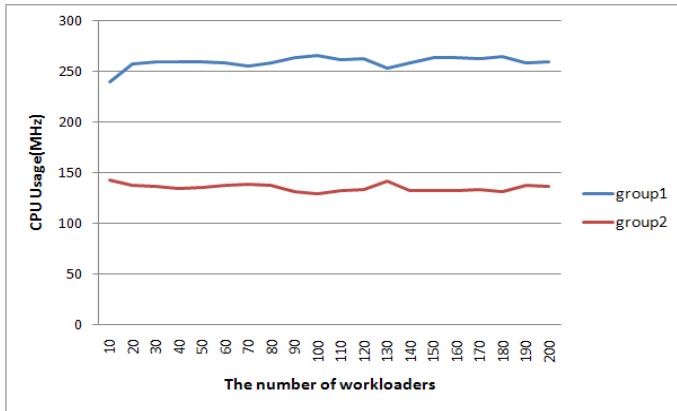


Fig. CPU usage in resource guaranteeing test

Group1 : group2 = 1(1) : 2(continuous increase up to 200)

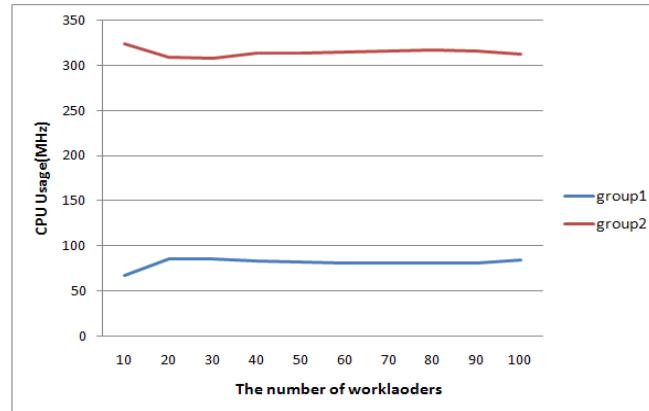


Fig. CPU usage in resource limitation test

Group1 : group2 : group3 : group4 : group5 = 1(100) : 1(100) : 1(100) : 1(100) : 1(100)

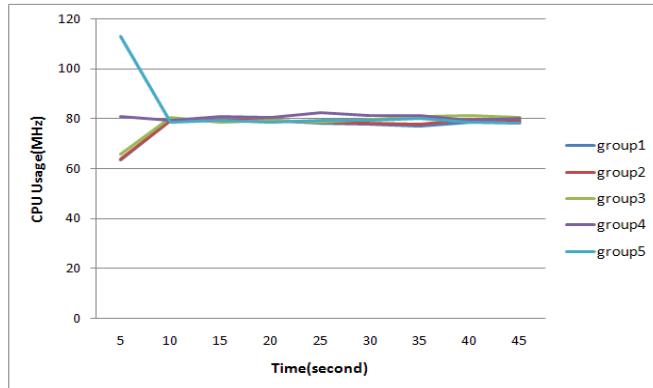


Fig. CPU usage in resource isolation test

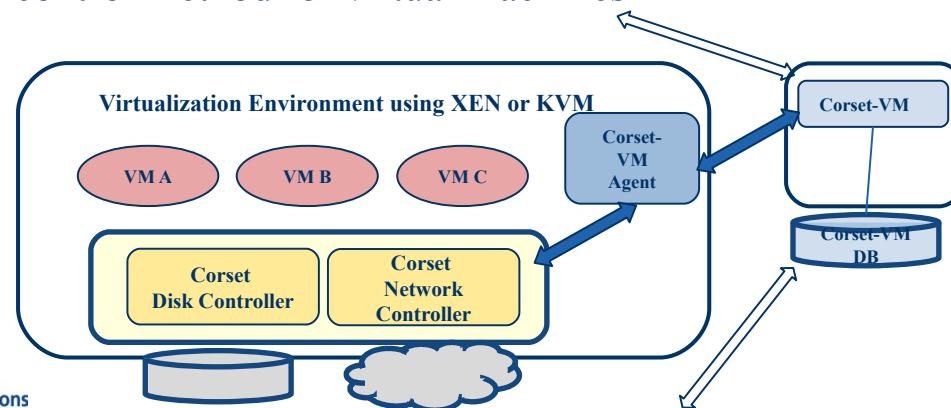
Concluding Remarks

❖ Features of Corset

- Suggested service-oriented resource management system
 - Allocates the system resources focused on service according to its priority
 - Supports the expectable resources to each service to maintain the its QoS
 - Proportional allocation
 - Fixed allocation
 - Range allocation
- Makes the limited system resources to be used efficiently
- Supports the dynamic controls of the system resources

❖ Future works

- Supports the I/O control method for virtual machines



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Thank You !

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